## ÉCOLE POLYTECHNIQUE FÉDÉRALE DE LAUSANNE

School of Computer and Communication Sciences

Handout 6
Binary Hypothesis testing

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## 1 Hypothesis testing

Consider the problem of deciding which of two hypotheses, hypothesis 0 or hypothesis 1, is true, based on an observation Y. For simplicity, assume that the observation Y is a random variable taking values in an alphabet  $\mathcal{Y}$  — a finite set of  $K = |\mathcal{Y}|$  letters — and under hypothesis j it has distribution  $P_j$ . To avoid trivial cases we will assume that for each  $y \in \mathcal{Y}$  both  $P_0(y)$  and  $P_1(y)$  are strictly positive. Otherwise, if we observe a y with, say,  $P_0(y) = 0$ , we would know for sure that hypothesis 1 is true.

A deterministic decision rule associates to each  $y \in \mathcal{Y}$  a binary value — i.e., the rule is a function  $\phi: \mathcal{Y} \to \{0,1\}$  — and we decide in favor of hypothesis  $\phi(y)$  if the observation Y equals y. In general, we will allow for randomized decision rules: such a rule is characterized by a function  $\phi: \mathcal{Y} \to [0,1]$  that associates to each  $y \in \mathcal{Y}$  a value in the *interval* [0,1], that gives the probability of deciding in favor of hypothesis 1. If our observation Y equals y, we flip a coin that comes heads with probability  $\phi(y)$  and tails with probability  $1 - \phi(y)$ , and decide accordingly: 1 if heads, 0 if tails. We will identify a decision rule with the function  $\phi$ .

In this set up there are two kinds of error: deciding 1 when the true hypothesis is 0, and deciding 0 when the true hypothesis is 1. Letting  $\pi_{\phi}(i|j)$  for rule  $\phi$  denote the probability of deciding i when the truth is j, we see that

$$\pi_{\phi}(0|1) = \sum_{y} P_1(y)[1 - \phi(y)], \quad \pi_{\phi}(1|0) = \sum_{y} P_0(y)\phi(y).$$

Given  $P_0$  and  $P_1$  and a positive real number  $\eta > 0$ , let  $\Phi_{\eta}$  to be the set of decision rules  $\phi$  of the form

$$\phi(y) = \begin{cases} 1 & \text{if } P_1(y) > \eta P_0(y) \\ 0 & \text{if } P_1(y) < \eta P_0(y). \end{cases}$$
 (1)

Note that if there is no y for which  $P_1(y) = \eta P_0(y)$ , the test  $\phi$  is uniquely specified and  $\Phi_{\eta}$  contains only this test.

LEMMA 1. The rules in  $\Phi_{\eta}$  are minimizers of  $\pi(0|1) + \eta \pi(1|0)$ .

*Proof.* For any rule  $\phi \in \Phi_n$ , as a consequence of (1), for every  $y \in \mathcal{Y}$ 

$$P_1(y)[1 - \phi(y)] + \eta P_0(y)\phi(y) = \min\{P_1(y), \eta P_0(y)\}.$$

Thus for any rule  $\phi \in \Phi_n$ 

$$\pi_{\phi}(0|1) + \eta \pi_{\phi}(1|0) = \sum_{y} P_{1}(y)[1 - \phi(y)] + \eta P_{0}(y)\phi(y) = \sum_{y} \min\{P_{1}(y), \eta P_{0}(y)\}.$$

Suppose now  $\psi$  is any decision rule. The lemma follows by noting that

$$\pi_{\psi}(0|1) + \eta \pi_{\psi}(1|0) = \sum_{y} P_1(y)[1 - \psi(y)] + \eta P_0(y)\psi(y) \ge \sum_{y} \min\{P_1(y), \eta P_0(y)\}. \quad \Box$$

THEOREM (NEYMAN-PEARSON, 1933). For any  $\alpha \in [0,1]$ , (i) there is a rule  $\phi$  of the form (1) such that  $\pi_{\phi}(1|0) = \alpha$ , and (ii) for any decision rule  $\psi$  either  $\pi_{\psi}(0|1) \geq \pi_{\phi}(0|1)$  or  $\pi_{\psi}(1|0) \geq \pi_{\phi}(1|0)$ .

Proof. Assertion (ii) follows from the lemma above: a  $\psi$  that violates both the inequalities would contradict the lemma. It thus suffices to prove (i), the existence of a rule  $\phi$  of the form (1) with  $\pi_{\phi}(1|0) = \alpha$ . To that end, define  $\Lambda(y) = P_1(y)/P_0(y)$ , and label the elements of  $\mathcal{Y}$  as  $\mathcal{Y} = \{y_1, \ldots, y_K\}$  such that  $\Lambda(y_1) \geq \Lambda(y_2) \geq \cdots \geq \Lambda(y_K)$ . Now define,  $a_i = \sum_{j=1}^i P_0(y_j)$  for  $i = 0, \ldots, K$ . We then have  $0 = a_0 < a_1 < \cdots < a_K = 1$ . Given  $0 \leq \alpha \leq 1$ , we can find  $1 \leq i \leq K$  for which  $a_{i-1} \leq \alpha \leq a_i$ , so that  $\alpha = (1 - \rho)a_{i-1} + \rho a_i$  for some  $\rho \in [0, 1]$ . Observe that  $\alpha = \sum_{j=1}^{i-1} P_0(y_j) + \rho P_0(y_i)$ , and that the rule

$$\phi(y) = \begin{cases} 1 & y \in \{y_1, \dots, y_{i-1}\} \\ \rho & y = y_i \\ 0 & y \in \{y_{i+1}, \dots, y_K\} \end{cases}$$

is of the form (1) with  $\eta = \Lambda(y_i)$ , and  $\pi_{\phi}(1|0) = \alpha$ .

Rules of the form (1) are based on a likelihood ratio test: they compare the likelihood ratio  $\Lambda(y) = P_1(y)/P_0(y)$  to a threshold  $\eta$  to make a decision. If the likelihood ratio is larger than the threshold, decide 1; if less, decide 0. Equivalently one may compare the log likelihood ratio,  $\log(P_1(y)/P_0(y))$  to the threshold  $\log \eta$ .

The theorem stated just above shows the dominant nature of likelihood ratio tests in making decisions: given any decision rule  $\psi$ , we can find a (log) likelihood ratio test  $\phi$  which is 'as good or better' — in the sense that the two error probabilities satisfy  $\pi_{\phi}(0|1) \leq \pi_{\psi}(0|1)$  and  $\pi_{\phi}(1|0) \leq \pi_{\psi}(1|0)$ .