PROBLEM 1.

(a) We have

$$\begin{aligned}\Pr[U(1) \neq U^n \mid U^n = u^n] &= \Pr[U(1) \neq u^n \mid U^n = u^n] \overset{(*)}{=} \Pr[U(1) \neq u^n] \\ &= 1 - \Pr[U(1) = u^n] = 1 - \prod_{i=1}^{n} \Pr[U(1)_i = u_i] = 1 - \prod_{i=1}^{n} p_U(u_i),\end{aligned}$$

where $(*)$ follows from the independence of $U(1)$ and U^n.

(b) An encoding failure happens if and only if $U(m) \neq U^n$ for every $m = 1, 2, \ldots, M$. Therefore,

$$\begin{aligned}\Pr[\text{“failure”} \mid U^n = u^n] &= \Pr[U(m) \neq U^n, \forall m = 1, \ldots, M \mid U^n = u^n] \\ &= \Pr[U(m) \neq u^n, \forall m = 1, \ldots, M \mid U^n = u^n] \\ &= \Pr[U(m) \neq u^n, \forall m = 1, \ldots, M] \\ &= \prod_{m=1}^{M} \left(1 - \prod_{i=1}^{n} p_U(u_i)\right) = \left(1 - \prod_{i=1}^{n} p_U(u_i)\right)^M\end{aligned}$$

(c) Note that if $u^n \in \mathcal{T}_\epsilon^n(p_U)$, then $\prod_{i=1}^{n} p_U(u_i) \geq 2^{-nH(U)(1+\epsilon)}$, which implies

$$\begin{aligned}\Pr[\text{“failure”} \mid U^n = u^n] &= \left(1 - \prod_{i=1}^{n} p_U(u_i)\right)^M \leq \left(1 - 2^{-nH(U)(1+\epsilon)}\right)^M \\ &\overset{(*)}{\leq} \exp\left(-M 2^{-nH(U)(1+\epsilon)}\right) = \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right).\end{aligned}$$

where $(*)$ follows from the hint. Therefore, we have

$$\begin{aligned}\Pr[\text{“failure”} \mid U^n \in \mathcal{T}_\epsilon^n(p_U)] &= \frac{\Pr[\text{“failure”}, U^n \in \mathcal{T}_\epsilon^n(p_U)]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &= \frac{\sum_{u^n \in \mathcal{T}_\epsilon^n(p_U)} \Pr[\text{“failure”}, U^n = u^n]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &= \frac{\sum_{u^n \in \mathcal{T}_\epsilon^n(p_U)} \Pr[\text{“failure”} \mid U^n = u^n] \Pr[U^n = u^n]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &\leq \frac{\sum_{u^n \in \mathcal{T}_\epsilon^n(p_U)} \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right) \Pr[U^n = u^n]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &= \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right) \frac{\sum_{u^n \in \mathcal{T}_\epsilon^n(p_U)} \Pr[U^n = u^n]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &= \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right) \frac{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]}{\Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)]} \\ &= \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right).\end{aligned}$$

(d) Assume $R > H(U)$, then there exists $\epsilon > 0$ such that $R > H(U)(1+\epsilon)$. We have

$$\begin{aligned}
\Pr[\text{“failure”}] &= \Pr[\text{“failure”}, U^n \in \mathcal{T}_\epsilon^n(p_U)] + \Pr[\text{“failure”}, U^n \notin \mathcal{T}_\epsilon^n(p_U)] \\
&= \Pr[\text{“failure”} \mid U^n \in \mathcal{T}_\epsilon^n(p_U)] \Pr[U^n \in \mathcal{T}_\epsilon^n(p_U)] + \Pr[\text{“failure”}, U^n \notin \mathcal{T}_\epsilon^n(p_U)] \\
&\leq \Pr[\text{“failure”} \mid U^n \in \mathcal{T}_\epsilon^n(p_U)] + \Pr[U^n \notin \mathcal{T}_\epsilon^n(p_U)] \\
&\leq \exp\left(-2^{nR - nH(U)(1+\epsilon)}\right) + \Pr[U^n \notin \mathcal{T}_\epsilon^n(p_U)].
\end{aligned}$$

Since $R > H(U)(1+\epsilon)$ both terms in the above go to 0 as $n \to \infty$. Hence, $\Pr[\text{“failure”}] \to 0$ as n gets large.

Problem 2. Let the input distribution be p. We thus have

$$p(-1) + p(0) + p(1) = 1 \qquad p(-1) \geq 0, p(0) \geq 0, p(1) \geq 0$$

(since p is a distribution) and, to satisfy $E[b(X)] \leq \beta$ we must have

$$p(-1) + p(1) = 1 - p(0) \leq \beta.$$

Moreover,

$$\begin{aligned}
I(X;Y) &= H(Y) - H(Y|X) \\
&\overset{(a)}{=} H(Y) - p(0) \\
&\overset{(b)}{\leq} 1 - p(0) \\
&\overset{(c)}{\leq} \min\{1, \beta\}.
\end{aligned}$$

where (a) follows because given $\{X = -1\}$ or $\{X = 1\}$ there is no uncertainity in Y while given $\{X = 0\}$, Y is uniformly distributed in $\{-1, 1\}$, (b) holds since Y is binary with equality if $p(-1) + \frac{1}{2}p(0) = p(1) + \frac{1}{2}p(0) = \frac{1}{2}$ (which happens if we choose $p(1) = p(-1) = \frac{1}{2}(1 - p(0))$) and (c) holds because of the cost constraint and is equality if we choose $p(0) = \max\{1 - \beta, 0\}$. Hence, the capacity is

$$C = \begin{cases} \beta, & \text{if } \beta \leq 1 \\ 1, & \text{if } \beta > 1 \end{cases}.$$

Problem 3.

(a) All rates less than $\frac{1}{2}\log_2(1 + \frac{P}{\sigma^2})$ are achievable.

(b) The new noise $Z_1 - \rho Z_2$ has zero mean and variance $\mathrm{E}((Z_1 - \rho Z_2)^2) = \sigma^2(1 - \rho^2)$. Therefore, all rates less than $\frac{1}{2}\log_2(1 + \frac{P}{\sigma^2(1-\rho^2)})$ are achievable.

(c) The capacity is $C = \max I(X; Y_1, Y_2) = \max(h(Y_1, Y_2) - h(Z_1, Z_2)) = \frac{1}{2}\log_2(1 + \frac{P}{\sigma^2(1-\rho^2)})$. This shows that the scheme used in (b) is a way to achieve capacity.

Problem 4.

(a) We have

$$I(X;Y) = h(Y) - h(Y|X) = h(Y) - h(Z|X) = h(Y) - h(Z).$$

where the last equality is because Z is independent of X.

(b) In the natural log basis,

$$h(Z) = -\int f_Z(z) \ln f_Z(z)\, dz = \int_0^\infty z e^{-z}\, dz = 1 \text{ nats.}$$

(c) Since $Y = X + Z$, the expectation of Y, $E[Y]$ equals $E[X] + E[Z]$. Since $E[X]$ is constrained to be less than or equal to P and $E[Z] = 1$, we see that $E[Y] \leq P + 1$. Since X is constrained to be non-negative and so is Z, we see that Y is also constrained to be non-negative.
From Homework 9, Problem 7 we know that among non-negative random variables of a given expectation λ, the one with density $p(y) = e^{-y/\lambda}/\lambda$ has the largest differential entropy. This differential entropy in natural units is

$$\int_0^\infty \frac{e^{-y/\lambda}}{\lambda} [\ln \lambda + y/\lambda]\, dy = \ln \lambda + 1 \text{ nats.}$$

Thus, the differential entropy of Y is less than $1 + \ln E[Y] \leq 1 + \ln(1 + P)$, which implies

$$C \leq \ln(1 + P) \text{ nats}$$

At this point, we do not know if Y can be made to have an exponential distribution with mean $1 + P$ so we cannot know if this above inequality is an equality or not.

(d) The Laplace transform of the random variable Y is $E(e^{sY}) = E(e^{s(X+Z)}) = E(e^{sX})E(e^{sZ})$, where the latter equality follows from the independence of X and Z. Therefore we have that $E(e^{sX}) = \frac{E(e^{sY})}{E(e^{sZ})}$. Computing $E(e^{sY})$,

$$\begin{aligned} E(e^{sY}) &= \int_\infty^\infty e^{sy} f_Y(y) dy \\ &= \int_0^\infty e^{sy} \mu e^{-\mu y} dy \\ &= \frac{\mu}{\mu - s} \quad \forall s \leq \mu \end{aligned}$$

The expectation is not defined for $s > \mu$ (as the integral blows up). Likewise, we evaluate $E(e^{sZ}) = \frac{1}{1-s}$ (defined for $s \leq 1$). Therefore for $s \leq \min(1, \mu)$, we can evaluate $E(e^{sX})$ as

$$\begin{aligned} E(e^{sX}) &= \frac{E(e^{sY})}{E(e^{sZ})} \\ &= \mu \frac{1 - s}{\mu - s} \\ &= \mu + (1 - \mu)\frac{\mu}{\mu - s} \end{aligned}$$

Inverting the Laplace transform $E(e^{sX})$ gives us the distribution of the X that gives an exponential distribution for Y. From inspection, we can deduce this distribution of X to be

$$f_X(x) = \mu\delta(x) + (1 - \mu)\mu e^{-\mu x} \quad x \geq 0$$

Notice that the distribution is a convex combination of the exponential distribution and the distribution that puts all the mass on one point (in this case the point $x = 0$).

(e) By taking $\mu = 1/(1+P)$, we see that there is a density on X which makes the density of Y an exponential with mean $1+P$. Furthermore, this density on X makes X non-negative, and, $E[X] = E[Y] - E[Z] = P$. Thus, the bound of part (c) can be achieved.

PROBLEM 5.

(a)

$$\begin{aligned} F(p, r_p) - F(p, r) &= \sum_{x\in\mathcal{X}} \sum_{y\in\mathcal{Y}} p(x)P(y|x) \log_2 \frac{r_p(x|y)}{r(x|y)} \\ &= \sum_{x\in\mathcal{X}} \sum_{y\in\mathcal{Y}} p(x)P(y|x) \log_2 \frac{p(x)P(y|x)}{r(x|y)\sum_{x'\in\mathcal{X}} p(x')P(y|x')} \\ &= D(P_1\|P_2) \geq 0, \end{aligned}$$

where $P_1(x,y) := p(x)P(y|x)$ and $P_2(x,y) := r(x|y)\sum_{x'\in\mathcal{X}} p(x')P(y|x')$.

(b) We can rewrite $F(p,r)$ as follows:

$$F(p,r) = \Big(\sum_{x\in\mathcal{X}}\sum_{y\in\mathcal{Y}} p(x)P(y|x)\log_2 r(x|y)\Big) + \Big(\sum_{x\in\mathcal{X}} p(x)\log_2\frac{1}{p(x)}\Big). \tag{1}$$

The first term in (1) is linear in p while the second term is strictly concave in p (since the function $t \longrightarrow t\log_2\frac{1}{t}$ is strictly concave). Therefore, $F(p,r)$ is strictly concave in p.

The first term in 1 is concave in r (since the function $\log_2$ is concave) and the second term is constant with respect to r. Therefore, $F(p,r)$ is concave in r.

(c) For every $x \in \mathcal{X}$, we have:

$$\frac{\partial F(p,r_k)}{\partial p(x)} = \sum_{y\in\mathcal{Y}} P(y|x)\log_2 r_k(x|y) + \log_2\frac{1}{p(x)} - \frac{1}{\ln 2}.$$

A probability distribution p satisfies the Kuhn-Tucker conditions if and only if there exists a real number λ such that for all $x \in \mathcal{X}$, we have $\dfrac{\partial F(p,r_k)}{\partial p(x)} \leq \lambda$ with equality if $p(x) > 0$. Therefore, for all $x \in \mathcal{X}$ we have:

$$\sum_{y\in\mathcal{Y}} P(y|x)\log_2 r_k(x|y) - \log_2(p(x)) \leq \lambda',$$

where $\lambda' = \lambda + \dfrac{1}{\ln 2}$. This shows that

$$p(x) \geq 2^{-\lambda'}\alpha_k(x).$$

If $p(x) > 0$, we have $p(x) = 2^{-\lambda'}\alpha_k(x)$, and if $p(x) = 0$ we must also have $p(x) = 2^{-\lambda'}\alpha_k(x) = 0$ since $2^{-\lambda'}2^{\sum_{y\in\mathcal{Y}} P(y|x)\log_2 r_k(x|y)} \geq 0$. We conclude that $p(x) = 2^{-\lambda'}\alpha_k(x)$ in all cases. Therefore, $1 = 2^{-\lambda'}\sum_{x\in\mathcal{X}}\alpha_k(x)$, and $\lambda' = \log_2\sum_{x\in\mathcal{X}}\alpha_k(x)$. We conclude that the only distribution that satisfies the Kuhn-Tucker conditions is the one given by $p(x) = \dfrac{\alpha_k(x)}{\sum_{x'\in\mathcal{X}}\alpha_k(x')}$. On the other hand, the fact that $F(p,r_k)$ is concave in p shows that it admits a maximum p_{k+1}, which has to satisfy the Kuhn-Tucker conditions. Therefore, $p_{k+1}(x) = \dfrac{\alpha_k(x)}{\sum_{x'\in\mathcal{X}}\alpha_k(x')}$.

(d) $C \geq F(p_{k+1}, r_{k+1})$ since $F(p_{k+1}, r_{k+1}) = I(X;Y)|_{p_X = p_{k+1}}$. This implies that $C \geq F(p_{k+1}, r_k)$ since $F(p_{k+1}, r_{k+1}) \geq F(p_{k+1}, r_k)$. On the other hand, we have

$$\begin{aligned}
&F(p_{k+1}, r_k) \\
&= \sum_{x\in\mathcal{X}} \sum_{y\in\mathcal{Y}} \frac{\alpha_k(x)}{\sum_{x'\in\mathcal{X}} \alpha_k(x')} P(y|x) \log_2 \frac{r_k(x|y) \sum_{x'\in\mathcal{X}} \alpha_k(x')}{\alpha_k(x)} \\
&= \sum_{x\in\mathcal{X}} \frac{\alpha_k(x)}{\sum_{x'\in\mathcal{X}} \alpha_k(x')} \Big[\sum_{y\in\mathcal{Y}} P(y|x) \log_2 r_k(x|y) - \log_2(\alpha_k(x)) + \log_2 \sum_{x'\in\mathcal{X}} \alpha_k(x') \Big] \\
&= \log_2 \sum_{x'\in\mathcal{X}} \alpha_k(x') + \sum_{x\in\mathcal{X}} \frac{\alpha_k(x)}{\sum_{x'\in\mathcal{X}} \alpha_k(x')} \Big[\log_2(\alpha_k(x)) - \log_2(\alpha_k(x)) \Big] \\
&= \log_2 \sum_{x\in\mathcal{X}} \alpha_k(x).
\end{aligned}$$

(e)

$$\begin{aligned}
\log_2 \frac{\alpha_k(x)}{p_k(x)} &= \log_2 \alpha_k(x) - \log_2 p_k(x) = \sum_{y\in\mathcal{Y}} P(y|x) \log_2 r_k(x|y) - \log_2 p_k(x) \\
&= \sum_{y\in\mathcal{Y}} P(y|x) \log_2 \frac{r_k(x|y)}{p_k(x)} = \sum_{y\in\mathcal{Y}} P(y|x) \log_2 \frac{P(y|x)}{\sum_{x'\in\mathcal{X}} p_k(x') P(y|x')}.
\end{aligned}$$

(f) Given that $\log_2 \dfrac{\alpha_k(x)}{p_k(x)} = \displaystyle\sum_{y\in\mathcal{Y}} P(y|x) \log_2 \frac{P(y|x)}{\sum_{x'\in\mathcal{X}} p_k(x') P(y|x')}$, the inequality $C \leq \displaystyle\sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{\alpha_k(x)}{p_k(x)}$ is a direct application of Homework 8 Problem 2.

(g) From (d) and (f), we have:

$$\begin{aligned}
&C - F(p_{k+1}, r_k) \\
&\leq \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{\alpha_k(x)}{p_k(x)} - \log_2 \sum_{x\in\mathcal{X}} \alpha_k(x) = \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{\alpha_k(x)}{p_k(x)} - \log_2 \sum_{x'\in\mathcal{X}} \alpha_k(x') \\
&= \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{\alpha_k(x)}{p_k(x) \sum_{x'\in\mathcal{X}} \alpha_k(x')} = \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{p_{k+1}(x)}{p_k(x)} \leq \max_{x\in\mathcal{X}} \log_2 \frac{p_{k+1}(x)}{p_k(x)}.
\end{aligned}$$

(h) We prove it by induction on n. The result is trivial for $n = 0$. Now assume that it is true for n, and let us prove it for $n + 1$:

$$\begin{aligned}
\sum_{k=0}^{n+1} (C - F(p_{k+1}, r_k)) &= C - F(p_{n+2}, r_{n+1}) + \sum_{k=0}^{n} (C - F(p_{k+1}, r_k)) \\
&\leq \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{p_{n+2}(x)}{p_{n+1}(x)} + \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{p_{n+1}(x)}{p_0(x)} \\
&= \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{p_{n+2}(x)}{p_0(x)}.
\end{aligned}$$

On the other hand, since $p_{n+1}(x) \leq 1$ for all $x \in \mathcal{X}$, we have:

$$\sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{p_{n+1}(x)}{p_0(x)} \leq \sum_{x\in\mathcal{X}} p^*(x) \log_2 \frac{1}{1/|\mathcal{X}|} = \log_2 |\mathcal{X}|.$$

(i) The sequence $s_n = \sum_{k=0}^{n} C - F(p_{k+1}, r_k)$ is increasing and upper-bounded, thus convergent, which implies that the sequence $C - F(p_{k+1}, r_k) = s_k - s_{k-1}$ converges to zero. Therefore, $F(p_{k+1}, r_k)$ converges to C.

Problem 6.

(a) Since C is non-empty, it contains some codeword x. By linearity C must contain $x + x$. But, for any x, $x + x$ is the all-zero sequence since we are doing modulo-2 sums. So, C contains the all-zero sequence.

(b) The elements of D' are those sequences of the form $x + y$ where y is in D. Since x is in C and D is a subset of C, any x and y are both in C, and so is their sum.

(c) Suppose there was an element z common to D and D'. Then $z = x + y$ where y is in D. Since we assumed that D is a linear subset, then $z + y$ is also in D. But $z + y$ equals x, and we arrive at the contradiction that x is in D.

(d) Since the mapping $y \mapsto x+y$ is a bijection, D and D' are in one-to-one correspondence, and hence have the same number of elements.

(e) Suppose z_1 and z_2 are in $D \cup D'$. There are four possibilities: (1) both z_1 and z_2 are in D, (2) both z_1 and z_2 are in D', (3) z_1 is in D, z_2 is in D', (4) z_1 is in D', z_2 is in D. In case (1), the linearity of D implies that $z_1 + z_2$ is in D. In case (2), $z_1 = x + y_1$ and $z_2 = x + y_2$ for some y_1 and y_2 both in D, then $z_1 + z_2 = x + x + y_1 + y_2 = y_1 + y_2$ is in D. In case (3) $z_2 = x + y_2$ and $z_1 + z_2 = x + (z_1 + y_2)$, which is in D', and similarly in case (4). Thus in all cases $z_1 + z_2$ is in $D \cup D'$ and we see that $D \cup D'$ is a linear subset of C.

(f) We thus see that if at the beginning of step (ii) D is a linear subset of C, at the end of step (iii) $D \cup D'$ is linear, is a subset of C because both D and D' are, and has twice as many elements of D since D' has the same number of elements of D and is disjoint from it. Thus, when the algorithm terminates, D contains all elements of C and since it is a subset of C it must equal C. Furthermore, its size, being equal to successive doublings of 1, is a power of 2.