ÉCOLE POLYTECHNIQUE FÉDÉRALE DE LAUSANNE
School of Computer and Communication Sciences

Handout 22 | Information Theory and Coding
Homework 9 | Nov. 21, 2017

PROBLEM 1. Let $f : U \to \mathbb{R}$ be a convex function on U and assume that there exists $a, b \in \mathbb{R}$ such that $a \leq f(x) \leq b$ for all $x \in U$. Let h be an increasing convex function defined on the interval $[a, b]$. Show that the function $g = h \circ f$ is convex on U.

PROBLEM 2. A function $f(v)$ is defined on a convex region R of a vector space. Show that $f(v)$ is convex iff the function $f(\lambda v_1 + (1 - \lambda)v_2)$ is a convex function of λ, $0 \leq \lambda \leq 1$, for all $v_1, v_2 \in R$.

PROBLEM 3. Show that among all non-negative random variables with mean λ the exponential random variable has the largest differential entropy. Hint: let $p(x) = e^{-x/\lambda}/\lambda$ be the density of the exponential random variable and let $q(x)$ be some other density with mean λ. Consider $D(q\|p)$ and mimic the proof in class for the maximal entropy of the Gaussian.

PROBLEM 4. Consider an additive noise channel with input $x \in \mathbb{R}$, and output

$$Y = x + Z$$

where Z is a real random variable independent of the input x, has zero mean and variance equal to σ^2.

In this problem we prove in two different ways that the Gaussian channel has the smallest capacity among all additive noise channels of a given noise variance. Let $\mathcal{N}_{\sigma^2}$ denote the Gaussian density with zero mean and variance σ^2.

FIRST METHOD

Let X be a Gaussian random variable with zero-mean and variance P. Let $\mathcal{N}_P$ denote its density $\mathcal{N}_P(x) = \frac{1}{\sqrt{2\pi P}} \mathrm{e}^{-\frac{x^2}{2P}}$.

(a) Show that $I(X; Y) = H(X) - H(X - \alpha Y|Y)$ for any $\alpha \in \mathbb{R}$.

(b) Show that $H(X - \alpha Y) \leq \frac{1}{2} \log 2\pi e E((X - \alpha Y)^2)$ for any $\alpha \in \mathbb{R}$.

(c) Deduce from (a) and (b) that

$$I(X; Y) \geq H(X) - \tfrac{1}{2} \log 2\pi e E((X - \alpha Y)^2)$$

for any $\alpha \in \mathbb{R}$.

(d) Show that $E((X - \alpha Y)^2) \geq \frac{\sigma^2 P}{\sigma^2 + P}$ with equality if and only if $\alpha = \frac{P}{P + \sigma^2}$.

(e) Deduce from (c) and (d) that

$$I(X; Y) \geq \tfrac{1}{2} \log \left(1 + \frac{P}{\sigma^2}\right)$$

and conclude that the Gaussian channel has the smallest capacity among all additive noise channels of a given noise variance.

SECOND METHOD

(a) Denote the input probability density by p_X. Verify that

$$I(X;Y) = \iint p_X(x) p_Z(y-x) \ln \frac{p_Z(y-x)}{p_Y(y)}\, dx dy \quad \text{nats.}$$

where p_Y is the density of the output when the input has density p_X.

(b) Now set $p_X = \mathcal{N}_P$. Verify that

$$\tfrac{1}{2} \ln\bigl(1 + P/\sigma^2\bigr) = \iint p_X(x) p_Z(y-x) \ln \frac{\mathcal{N}_{\sigma^2}(y-x)}{\mathcal{N}_{P+\sigma^2}(y)}\, dx dy.$$

(c) Still with $p_X = \mathcal{N}_P$, show that

$$\tfrac{1}{2} \ln(1 + P/\sigma^2) - I(X;Y) \le 0.$$

[Hint: use (a) and (b) and $\ln t \le t - 1$.]

(d) Show that an additive noise channel with noise variance σ^2 and input power P has capacity at least $\frac{1}{2}\log_2(1 + P/\sigma^2)$ bits per channel use. Conclude that the Gaussian channel has the smallest capacity among all additive noise channels of a given noise variance.