## Random matrices and communication systems: WEEK 14

## 1 Diversity-multiplexing tradeoff (cont'd)

Remember from last lecture that we are after computing the diversity order of a multiple antenna channel:

$$d(r) = \lim_{P \to \infty} -\frac{\log(P_{\text{out}}(r \log P))}{\log P} = \lim_{P \to \infty} -\frac{\log(\mathbb{P}(\log \det(I + PHH^*) < r \log P)))}{\log P}$$

where H is an  $n \times n$  matrix with i.i.d.  $\sim \mathcal{N}_{\mathbb{C}}(0,1)$  entries (and n is fixed). The above probability can be rewritten as

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) = \mathbb{P}\left(\sum_{j=1}^n \log(1 + P\lambda_j) < r \log P\right)$$

where  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of the matrix  $HH^*$ , which are all non-negative. In Lecture 6, we have seen that the joint distribution of these eigenvalues is given by

$$p(\lambda_1, \dots, \lambda_n) = c_n \prod_{j=1}^n e^{-\lambda_j} \prod_{j < k} (\lambda_k - \lambda_j)^2 \text{ for } \lambda_1, \dots, \lambda_n \ge 0$$

where  $c_n$  is some positive constant. Using this, the above probability can be further rewritten as an n-fold integral:

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) = \int_{D_{\lambda}(r)} p(\lambda_1, \dots, \lambda_n) \, d\lambda_1 \cdots d\lambda_n$$

where

$$D_{\lambda}(r) = \left\{ 0 \le \lambda_1 \le \ldots \le \lambda_n : \sum_{j=1}^n \log(1 + P\lambda_j) < r \log P \right\}$$

(notice that the eigenvalues  $\lambda_1, \ldots, \lambda_n$  are ordered in increasing order here). The explicit computation of this integral remains a challenge, because of the highly correlated nature of the eigenvalues. We will see below that taking the high SNR limit  $(P \to \infty)$  in the above expression allows to drastically simplify the analysis. To this end, let us make the change of variables:

$$\lambda_j = P^{-\alpha_j} = \exp(-\alpha_j \log P), \text{ so } d\lambda_j = -(\log P) \exp(-\alpha_j \log P) d\alpha_j$$

This change of variable, even though depending on P, is perfectly valid for given value of P, and therefore also in the limit  $P \to \infty$  (provided some care is taken here). This gives rise to the following expression for the above probability:

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) = \int_{D_{\alpha}(r)} q(\alpha_1, \dots, \alpha_n) \, d\alpha_1 \cdots d\alpha_n$$

where

$$q(\alpha_1, \dots, \alpha_n) = c_n \exp\left(-\sum_{j=1}^n P^{-\alpha_j}\right) \prod_{j < k} \left(P^{-\alpha_j} - P^{-\alpha_k}\right)^2 \left(\log P\right)^n \exp\left(-\sum_{j=1}^n \alpha_j \log P\right)$$

and

$$D_{\alpha}(r) = \left\{ \alpha_1 \ge \ldots \ge \alpha_n \ (\alpha_j \in \mathbb{R}) \ : \ \sum_{j=1}^n \log\left(1 + P^{1-\alpha_j}\right) < r \ \log P \right\}$$

So far, these are exact expressions. We will now make a series of approximations which are valid in the limit  $P \to \infty$  (and which can all be rigorously justified by taking upper and lower bounds).

First observe that

$$\exp\left(-P^{-\alpha_j}\right) \begin{cases} \text{decays super-polynomially to zero} & \text{if } \alpha_j < 0\\ \text{tends to } 1 & \text{if } \alpha_j \ge 0 \end{cases}$$

so we may restrict the domain of integration  $D_{\alpha}(r)$  to its positive part where  $\alpha_1 \ge \ldots \ge \alpha_n \ge 0$ . Next, observe that

$$\log\left(1+P^{1-\alpha_j}\right) \simeq \begin{cases} (1-\alpha_j) \log P & \text{if } \alpha_j \leq 1\\ 0 & \text{if } \alpha_j > 1 \end{cases}$$

so  $\log(1+P^{1-\alpha_j})\simeq(1-\alpha_j)^+\log P$ . We can therefore replace the domain of integration  $D_{\alpha}(r)$  by

$$\widetilde{D}_{\alpha}(r) = \left\{ \alpha_1 \ge \ldots \ge \alpha_n \ge 0 : \sum_{j=1}^n (1 - \alpha_j)^+ \le r \right\}$$

and the above probability can be rewritten as

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) \doteq \int_{\widetilde{D}_{\alpha}(r)} \widetilde{q}(\alpha_1, \dots, \alpha_n) \, d\alpha_1 \cdots d\alpha_n$$

where

$$\widetilde{q}(\alpha_1, \dots, \alpha_n) = c_n \prod_{j < k} \left( P^{-\alpha_j} - P^{-\alpha_k} \right)^2 \left( \log P \right)^n \exp\left( -\sum_{j=1}^n \alpha_j \log P \right)$$

Furthermore, let us notice that  $c_n(\log P)^n \doteq 1$ , as  $\lim_{P\to\infty} \frac{\log(c_n(\log P)^n)}{\log P} = 0$ . Here comes now the "magic" trick: for  $\alpha_1 > \ldots > \alpha_n$ , we have

$$\prod_{j < k} \left( P^{-\alpha_j} - P^{-\alpha_k} \right)^2 \doteq \prod_{j < k} P^{-2\alpha_k} = \prod_{k=1}^n P^{-2(k-1)\alpha_k} = \exp\left( -\sum_{k=1}^n 2(k-1)\alpha_k \log P \right)$$

This implies that

$$\widetilde{q}(\alpha_1,\ldots,\alpha_n) \doteq \exp\left(-\sum_{j=1}^n (2j-1)\,\alpha_j\,\log P\right)$$

In this expression, we see that in the limit  $P \to \infty$ , the exponents  $\alpha_j$  become so to speak "independent". Finally, we obtain

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) \doteq \int_{\widetilde{D}_{\alpha}(r)} \exp\left(-\sum_{j=1}^n (2j-1) \alpha_j \log P\right) d\alpha_1 \cdots d\alpha_n$$

This expression can in turn be rewritten as

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) \doteq \int_{\widetilde{D}_{\alpha}(r)} P^{-f(\alpha)} d\alpha_1 \cdots d\alpha_n$$

where

$$f(\alpha) = \sum_{j=1}^{n} (2j-1) \alpha_j$$

Using then Laplace's integration method, we obtain

$$\mathbb{P}(\log \det(I + PHH^*) < r \log P) \doteq P^{-d(r)}$$

where the diversity order d(r) is given by

$$d(r) = \inf_{\widetilde{D}_{\alpha}(r)} f(\alpha) = \inf_{\substack{\alpha_1 \ge \dots \ge \alpha_n \ge 0\\ \sum_{j=1}^n (1-\alpha_j)^+ < r}} \sum_{j=1}^n (2j-1) \alpha_j$$

Doing this, we have therefore transformed the initial problem of evaluating an *n*-fold integral (in the limit  $P \to \infty$ ) into a simple linear optimization problem.

For n = 2, the problem reads

$$d(r) = \inf_{\substack{\alpha_1 \ge \alpha_2 \ge 0\\(1-\alpha_1)^+ + (1-\alpha_2)^+ < r}} \alpha_1 + 3 \alpha_2$$

whose solution is given by

$$\begin{cases} 0 \le r \le 1 : & \alpha_1 = 1, \, \alpha_2 = 1 - r, \, d(r) = 4 - 3r \\ 1 \le r \le 2 : & \alpha_1 = 2 - r, \, \alpha_2 = 0, \, d(r) = 2 - r \end{cases}$$

For low multiplexing gain  $(0 \le r \le 1)$ , outage occurs when both eigenvalues  $\lambda_1, \lambda_2$  of  $HH^*$  are small (more precisely,  $\lambda_1 \simeq P^{-1}$  and  $\lambda_2 \simeq P^{r-1}$ ), while for higher multiplexing gain  $(1 \le r \le 2)$ , outage occurs when only the smallest eigenvalue  $\lambda_2$  is small (more precisely  $\lambda_1 \simeq 1$  and  $\lambda_2 \simeq P^{r-2}$ ). As expected, the diversity drops to zero for values of r larger than or equal to 2 (as in this case, the target rate is higher than the ergodic capacity). On the figure below, the diversity order is drawn as a function of the multiplexing gain r, which illustrates the tradefoff between diversity and multiplexing.



For general values of n, the curve d(r) is the piecewise linear curve such that  $d(k) = (n-k)^2$  at integer values of r (so  $d(0) = n^2$  and d(n) = 0). Notice that the maximum diversity  $d = n^2$  corresponding to r = 0 matches the number of independent random variables in the channel matrix H.