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School of Computer and Communication Sciences

Principles of Digital Communications:
Summer Semester 2012

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Solution of Homework 5

Problem 1. (*QAM with Erasure*)

$$\begin{aligned} P_{00} &= Pr(\{N_1 \geq -a\} \cap \{N_2 \geq -a\}) \\ &= Pr(\{N_1 \leq a\})Pr(\{N_2 \leq a\}) \\ &= \left[1 - Q\left(\frac{a}{\sigma}\right)\right]^2. \end{aligned}$$

By symmetry:

$$\begin{aligned} P_{01} = P_{03} &= Pr(\{N_1 \leq -(2b - a)\} \cap \{N_2 \geq -a\}) \\ &= Pr(\{N_1 \geq 2b - a\})Pr(\{N_2 \leq a\}) \\ &= Q\left(\frac{2b - a}{\sigma}\right) \left[1 - Q\left(\frac{a}{\sigma}\right)\right]. \end{aligned}$$

$$\begin{aligned} P_{04} &= Pr(\{N_1 \leq -(2b - a)\} \cap \{N_2 \leq -(2b - a)\}) \\ &= Pr(\{N_1 \geq 2b - a\} \cap \{N_2 \geq 2b - a\}) \\ &= \left[Q\left(\frac{2b - a}{\sigma}\right)\right]^2. \end{aligned}$$

$$\begin{aligned} P_{0\delta} &= 1 - Pr(\{Y \in \mathcal{R}_0\} \cup \{Y \in \mathcal{R}_1\} \cup \{Y \in \mathcal{R}_2\} \cup \{Y \in \mathcal{R}_3\} | \mathbf{s}_0 \text{ was sent}) \\ &= 1 - P_{00} - P_{01} - P_{02} - P_{03} \\ &= 1 - \left[1 - Q\left(\frac{a}{\sigma}\right)\right]^2 - 2Q\left(\frac{2b - a}{\sigma}\right) \left[1 - Q\left(\frac{a}{\sigma}\right)\right] - \left[Q\left(\frac{2b - a}{\sigma}\right)\right]^2 \\ &= 1 - \left[1 - Q\left(\frac{a}{\sigma}\right) + Q\left(\frac{2b - a}{\sigma}\right)\right]^2 \end{aligned}$$

Equivalently,

$$\begin{aligned}
P_{0\delta} &= Pr(\{N_1 \in [a, 2b - a]\} \cup \{N_2 \in [a, 2b - a]\}) \\
&= Pr(N_1 \in [a, 2b - a]) + Pr(N_2 \in [a, 2b - a]) - Pr(\{N_1 \in [a, 2b - a]\} \cap \{N_2 \in [a, 2b - a]\}) \\
&= 2 \left[Q\left(\frac{2b - a}{\sigma}\right) - Q\left(\frac{a}{\sigma}\right) \right] - \left[Q\left(\frac{2b - a}{\sigma}\right) - Q\left(\frac{a}{\sigma}\right) \right]^2,
\end{aligned}$$

which gives the same result as before.

Problem 2. (*Gaussian Hypothesis testing*)

1. (a)

$$\begin{aligned}
L(Y_1, \dots, Y_n) &= \frac{P(Y_1, \dots, Y_n | H_1)}{P(Y_1, \dots, Y_n | H_0)} \\
&= e^{\frac{|Y - \mu_0|^2 - |Y - \mu_1|^2}{2\sigma^2}} \\
&= e^{\frac{2(\mu_1 - \mu_0)^T Y + \|\mu_0\|^2 - \|\mu_1\|^2}{2\sigma^2}}.
\end{aligned}$$

Taking the logarithm on both sides we obtain the simplified decision rule

$$(\mu_1 - \mu_0)^T Y \underset{H_1}{\overset{H_0}{\gtrless}} \frac{\|\mu_1\|^2 - \|\mu_0\|^2}{2}$$

If we denote $a = \mu_1 - \mu_0$ and $b = \|\mu_1\|^2 - \|\mu_0\|^2$ we see that $a^T Y = b$ characterizes an n -dimensional hyperplane which passes through $\frac{\mu_1 + \mu_0}{2}$ and separates the two decision regions.

(b) Under H_1 ,

$$(\mu_1 - \mu_0)^T Y \sim N((\mu_1 - \mu_0)^T \mu_1, \sigma^2 \|\mu_1 - \mu_0\|^2)$$

Hence we obtain

$$\begin{aligned}
P(E | H_1) &= Q\left(\frac{(\mu_1 - \mu_0)^T \mu_1 - \left(\frac{\|\mu_1\|^2 - \|\mu_0\|^2}{2}\right)}{\sqrt{\sigma^2 \|\mu_1 - \mu_0\|^2}}\right) \\
&= Q\left(\frac{\|\mu_1 - \mu_0\|}{2\sigma}\right) \\
&= Q\left(\frac{d}{2\sigma}\right),
\end{aligned}$$

where $d = \|\mu_1 - \mu_0\|$ is the distance between the mean vectors. Similarly we can show that $P(E | H_0) = Q\left(\frac{d}{2\sigma}\right)$. Hence, $P(E) = Q\left(\frac{d}{2\sigma}\right)$.

2. (a)

$$\begin{aligned} L(Y_1, \dots, Y_n) &= \frac{P(Y_1, \dots, Y_n | H_1)}{P(Y_1, \dots, Y_n | H_0)} \\ &= \left(\frac{\sigma_0}{\sigma_1} \right)^n e^{\left(\frac{\sigma_1^2 - \sigma_0^2}{2\sigma_1^2\sigma_0^2} \sum_{i=1}^n Y_i^2 \right)} \underset{H_1}{\overset{H_0}{\leq}} 1. \end{aligned}$$

Taking logarithm and using $\sigma_1 > \sigma_0$ we obtain the simplified decision rule

$$\sum_{i=1}^n y_i^2 \underset{H_1}{\overset{H_0}{\leq}} \frac{2\sigma_1^2\sigma_0^2}{\sigma_1^2 - \sigma_0^2} \left(-n \log \frac{\sigma_0}{\sigma_1} \right) = \frac{2n\sigma_1^2\sigma_0^2}{\sigma_1^2 - \sigma_0^2} \log \left(\frac{\sigma_1}{\sigma_0} \right).$$

The decision boundary is the n-dimensional hyper sphere where the points inside the sphere belong to H_0 whereas the points outside belong to H_1 .

(b) i. Let define $Z = Y_1^2 + Y_2^2$. We first obtain the CDF of Z as follows

$$\begin{aligned} F_Z(z) &= P(Y_1^2 + Y_2^2 \leq z) \\ &= \int_{Y_1^2 + Y_2^2 \leq z} \frac{1}{2\pi\sigma^2} e^{-\frac{y_1^2 + y_2^2}{2\sigma^2}} dy_1 dy_2 \end{aligned}$$

And we apply the following change of variables: $y_1 = r \sin(\theta)$ and $y_2 = r \cos(\theta)$ (thus, $r^2 = y_1^2 + y_2^2$) and obtain:

$$\begin{aligned} &= \int_{r=0}^{\sqrt{z}} \int_{\theta=0}^{2\pi} \frac{1}{2\pi\sigma^2} e^{-\frac{r^2}{2\sigma^2}} |det J(\theta, r)| dr d\theta \\ &= \int_{r=0}^{\sqrt{z}} \frac{1}{2\pi\sigma^2} e^{-\frac{r^2}{2\sigma^2}} 2\pi r dr \\ &= \left[-e^{-\frac{r^2}{2\sigma^2}} \right]_0^{\sqrt{z}} \\ &= 1 - e^{-\frac{z}{2\sigma^2}} \end{aligned}$$

Taking the derivative we obtain the probability density function $f_Z(z) = \frac{1}{2\sigma^2} e^{-\frac{z}{2\sigma^2}}$, which is an exponential with parameter $\frac{1}{2\sigma^2}$.

ii. For $n = 2$ and H_1 , we have

$$\begin{aligned} P(E|H_1) &= P\left(Y_1^2 + Y_2^2 \geq \frac{4\sigma_1^2\sigma_0^2}{\sigma_1^2 - \sigma_0^2} \log\left(\frac{\sigma_1}{\sigma_0}\right) | H_1 \right) \\ &= \int_0^{\frac{4\sigma_1^2\sigma_0^2}{\sigma_1^2 - \sigma_0^2} \log\left(\frac{\sigma_1}{\sigma_0}\right)} \frac{1}{2\sigma_1^2} e^{-\frac{z}{2\sigma_1^2}} dz \\ &= 1 - e^{-\frac{4\sigma_0^2}{\sigma_1^2 - \sigma_0^2} \log\left(\frac{\sigma_1}{\sigma_0}\right)} = 1 - \left(\frac{\sigma_0}{\sigma_1} \right)^{\frac{4\sigma_0^2}{\sigma_1^2 - \sigma_0^2}} \end{aligned}$$

Similarly, we obtain

$$\begin{aligned} P(E|H_0) &= \int_{\frac{4\sigma_1^2\sigma_0^2}{\sigma_1^2-\sigma_0^2} \log \frac{\sigma_1}{\sigma_0}}^{\infty} \frac{1}{2\sigma_0^2} e^{-\frac{z}{2\sigma_0^2}} dz \\ &= e^{-\frac{4\sigma_1^2}{\sigma_1^2-\sigma_0^2} \log \frac{\sigma_1}{\sigma_0}} = \left(\frac{\sigma_0}{\sigma_1}\right)^{\frac{4\sigma_1^2}{\sigma_1^2-\sigma_0^2}} \end{aligned}$$

Let us denote $\rho = \frac{\sigma_1}{\sigma_0}$. Then we have :

$$\begin{aligned} P(E|H_0) &= 1 - \left(\frac{1}{\rho}\right)^{\frac{4}{\rho^2-1}}, \quad P(E|H_1) = \left(\frac{1}{\rho}\right)^{\frac{4\rho^2}{\rho^2-1}} \\ \text{and } P(E) &= \frac{1}{2} + \frac{1}{2} \left(\left(\frac{1}{\rho}\right)^{\frac{4\rho^2}{\rho^2-1}} - \left(\frac{1}{\rho}\right)^{\frac{4}{\rho^2-1}} \right) \end{aligned}$$

iii. Let us denote $A(\rho) = \left(\frac{1}{\rho}\right)^{\frac{4\rho^2}{\rho^2-1}}$ and $B(\rho) = \left(\frac{1}{\rho}\right)^{\frac{4}{\rho^2-1}}$. Using the Hopital rule we obtain:

$$\lim_{\rho \rightarrow \infty} \log B(\rho) = \lim_{\rho \rightarrow \infty} \frac{-4 \log(\rho)}{\rho^2 - 1} = 0 \Rightarrow \lim_{\rho \rightarrow \infty} B(\rho) = 1, \text{ and } \lim_{\rho \rightarrow \infty} A(\rho) = 0^4 = 0$$

Hence

$$\lim_{\rho \rightarrow \infty} P(E) = 0.$$

Problem 3. (*Repeat Codes and Bhattacharyya Bound*)

1. First, we find the probability mass function of (W_1, \dots, W_N) given each of the two hypotheses:

$$\begin{aligned} p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 0) &= Pr\{\text{sgn}(X_1 + Z_1) = w_1, \dots, \text{sgn}(X_N + Z_N) = w_N | H = 0\} \\ &= Pr\{\text{sgn}(X_1 + Z_1) = w_1, \dots, \text{sgn}(X_N + Z_N) = w_N \mid \\ &\quad (X_1, \dots, X_N) = (1, \dots, 1)\} \quad (1) \\ &= Pr\{\text{sgn}(1 + Z_1) = w_1, \dots, \text{sgn}(1 + Z_N) = w_N\} \quad (2) \\ &= Pr\{\text{sgn}(1 + Z_1) = w_1\} \cdot \dots \cdot Pr\{\text{sgn}(1 + Z_N) = w_N\}. \quad (3) \end{aligned}$$

But since the Z_i are independent of each other and have the same distribution (or, as we say more frequently, since the Z_i are iid random variables), we can write this also as

$$p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 0) = \prod_{i=1}^N Pr\{\text{sgn}(1 + Z) = w_i\}. \quad (4)$$

Notice that the event $\{W_i = \text{sgn}(1 + Z) = 0\}$ has probability zero. Therefore, it is of no interest to our consideration. This means that the random variables W_i can only assume values 1 or -1 . Suppose that (w_1, \dots, w_N) contains k values of 1, and thus $(N - k)$ values of -1 . With this definition, we can rewrite

$$p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 0) = (\Pr\{1 + Z \geq 0\})^k (\Pr\{1 + Z \leq 0\})^{N-k}. \quad (5)$$

Let us introduce the following notation:

$$\epsilon \stackrel{\text{def}}{=} \Pr\{(1 + Z) \leq 0\} = 1 - Q\left(-\frac{1}{\sigma}\right) = Q\left(\frac{1}{\sigma}\right). \quad (6)$$

Then, we can write

$$p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 0) = (1 - \epsilon)^k \epsilon^{N-k}. \quad (7)$$

Under hypothesis $H = 1$, we have essentially the same derivation. Let us give only a few steps, using the same definitions of k (number of ones) and ϵ as above:

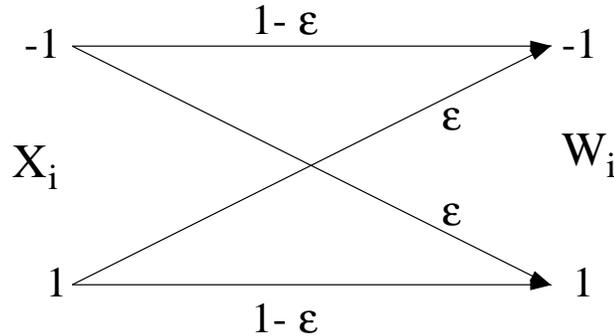
$$p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 1) = \prod_{i=1}^N \Pr\{\text{sign}(-1 + Z) = w_i\} \quad (8)$$

$$= (\Pr\{-1 + Z \geq 0\})^k (\Pr\{-1 + Z \leq 0\})^{N-k} \quad (9)$$

$$= \epsilon^k (1 - \epsilon)^{N-k} \quad (10)$$

Thus, we see that k is a sufficient statistic.

At this point, we give a simple picture that allows to derive all of this much more easily: The overall system between X and W may be viewed as a channel with input 1 or -1 and output also 1 or -1 . There is a certain probability ϵ (called *transition* or *crossover* probability) that the channel converts 1 into -1 or vice versa:



This particular channel is called the *Binary Symmetric Channel*. From this figure, various results can be found easily. For instance, it is clear that if we put N consecutive values 1 onto the channel, the probability of getting, at the output, a particular sequence (w_1, \dots, w_N) which contains exactly k values of 1 is simply $(1 - \epsilon)^k \epsilon^{N-k}$.

Similarly, the probability of getting, at the output, *any* sequence that contains exactly k values of 1 is $\binom{N}{k}(1-\epsilon)^k\epsilon^{N-k}$ because there are $\binom{N}{k}$ distinct particular sequences with exactly k ones each, and every one of them has probability $(1-\epsilon)^k\epsilon^{N-k}$.

Next, we have to develop the likelihood ratio:

$$L(w_1, \dots, w_N) = \frac{p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 1)}{p_{W_1 \dots W_N | H}(w_1, \dots, w_N | 0)} \quad (11)$$

$$= \frac{\epsilon^k (1-\epsilon)^{N-k}}{(1-\epsilon)^k \epsilon^{N-k}} = \left(\frac{\epsilon}{1-\epsilon} \right)^{2k-N}. \quad (12)$$

In terms of the loglikelihood ratio, we find

$$\log L(w_1, \dots, w_N) = (2k - N) \log \left(\frac{\epsilon}{1-\epsilon} \right) \stackrel{1}{\underset{0}{\gtrless}} 0. \quad (13)$$

Since $\epsilon < 1/2$, we know that $\log \left(\frac{\epsilon}{1-\epsilon} \right) < 0$, and thus, when we divide by this term, the direction of the inequality is changed. Using this, the decision rule can be written as

$$k \stackrel{1}{\underset{0}{\gtrless}} \frac{N}{2}. \quad (14)$$

That is, the best decision rule is simply *majority voting*: if the majority of the received values is 1, we decide for hypothesis $H = 0$ (i.e. transmitted value was 1); on the other hand, if the majority of the received values is -1 , we decide for hypothesis $H = 1$ (i.e. transmitted value was -1).

2. Let us assume that N is odd. Then,

$$P_e(0) = Pr\{ \text{there are less than } N/2 \text{ ones in the received sequence} \mid N \text{ ones were transmitted} \} \quad (15)$$

$$= \sum_{m=0}^{(N-1)/2} \binom{N}{m} (1-\epsilon)^m \epsilon^{N-m} \quad (16)$$

By the symmetry of the problem, $P_e(1)$ turns out to be the exact same expression, thus

$$P_e = \sum_{m=0}^{(N-1)/2} \binom{N}{m} (1-\epsilon)^m \epsilon^{N-m} \quad (17)$$

In case N is even, we introduce a slight asymmetry because the term for $N/2$ has to be assigned to either $H = 0$ or $H = 1$ (cannot be assigned to both).

Clearly, this sum cannot be evaluated explicitly. There are various techniques to bound it. In this homework, we consider the *Bhattacharyya bound* as encountered in class.

3. In class, you have seen the following derivation of the Bhattacharyya bound. For the optimal decision rule, we can write the probability of error as follows:

$$P_e = \sum_w \min_w \{p_{W|H}(w|0)p_H(0), p_{W|H}(w|1)p_H(1)\}, \quad (18)$$

where the sum is over all possible sequences w of length N . In our case, since $w_i \in \{-1, 1\}$, we have $w \in \{-1, 1\}^N$, and thus there are 2^N terms in the sum. But since for $a, b \geq 0$, we have that $\min a, b \leq \sqrt{ab}$, we get the following simple upper bound:

$$Pr\{e\} \leq \sum_w \sqrt{p_{W|H}(w|0)p_H(0)p_{W|H}(w|1)p_H(1)} \quad (19)$$

$$= \sqrt{p_H(0)p_H(1)} \sum_w \sqrt{p_{W|H}(w|0)p_{W|H}(w|1)} \quad (20)$$

$$\leq \frac{1}{2} \sum_w \sqrt{p_{W|H}(w|0)p_{W|H}(w|1)}, \quad (21)$$

where the last inequality follows because for $c, d \geq 0$, we have $\sqrt{cd} \leq (c + d)/2$. Now we have to plug in:

$$\tilde{P}_e \leq \frac{1}{2} \sum_w \sqrt{p_{W|H}(w|0)p_{W|H}(w|1)} \quad (22)$$

$$= \frac{1}{2} \sum_w \sqrt{(1 - \epsilon)^{k(w)} \epsilon^{N-k(w)} \epsilon^{k(w)} (1 - \epsilon)^{N-k(w)}}, \quad (23)$$

where we used $k(w)$ to denote the number of values 1 in the sequence w . We find furthermore

$$\tilde{P}_e \leq \frac{1}{2} \sum_w \sqrt{\epsilon^N (1 - \epsilon)^N} = \frac{1}{2} \sqrt{\epsilon^N (1 - \epsilon)^N} \sum_w 1 \quad (24)$$

$$= \frac{1}{2} \sqrt{\epsilon^N (1 - \epsilon)^N} 2^N = \frac{1}{2} \left(2\sqrt{\epsilon(1 - \epsilon)}\right)^N. \quad (25)$$

4. Again, we assume that N is odd; note however that the case when N is even would not add much insight. We used the following matlab program:

```
%
% Principles of Digital Communications, Summer Semester 2001 (Prof. B. Rimoldi)
%
% Solution to Homework 4, Problem 1.(iv)
% Michael Gastpar
%
% Notation: we are using Pe1 for $P_e$ and Pe2 for $\tilde{P}_e$
%
N = [ 1:2:30 ];
sigma = 1;
```

```

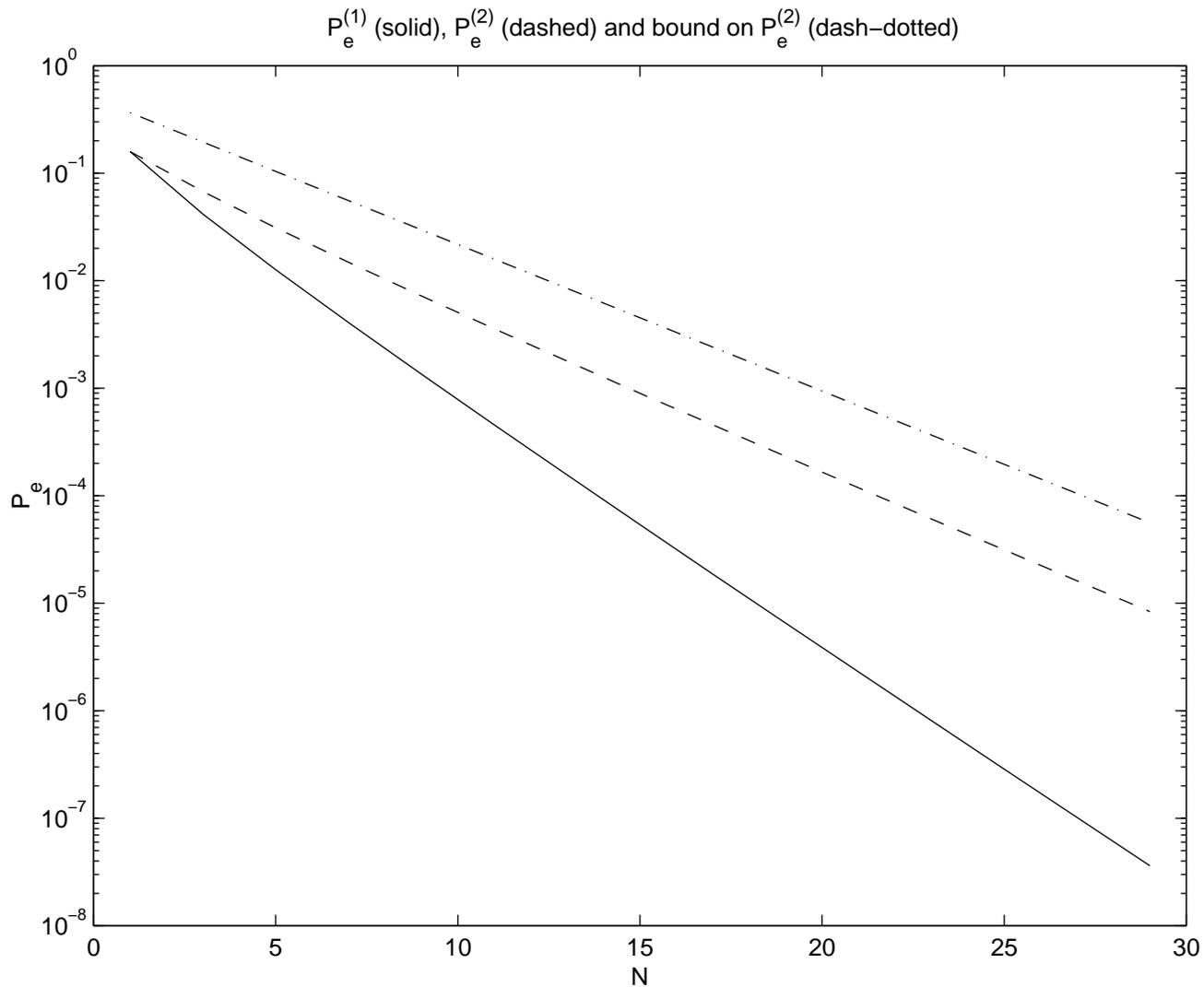
Pe1 = 1/2 * erfc( sqrt(N)/sigma /sqrt(2))

epsilon = 1/2 * erfc( 1/sigma /sqrt(2));
Pe2 = zeros(1, length(N));
for ic = 1:length(N),
    for m = 0:(N(ic)-1)/2,
        Pe2(ic) = Pe2(ic) + prod(N(ic)-m+1:N(ic))/prod(1:m) * (1-epsilon)^m * epsilon^(N(ic)-m);
    end;
end;
Pe2Bhatt = 1/2 * (2*sqrt(epsilon*(1-epsilon))).^N

semilogy( N, Pe1, N, Pe2Bhatt, '--', N, Pe2, '-.');
title('P_e^{(1)} (solid), P_e^{(2)} (dashed) and bound on P_e^{(2)} (dash-dotted)');
xlabel('N');
ylabel('P_e');

```

which gives the following plot for the error probabilities:



Problem 4. (*Tighter Union Bhattacharyya Bound: Binary Case*)

1. From the definition of the decision region \mathcal{R}_i

$$\mathcal{R}_i = \{y : P_H(i)f_{Y|H}(y|i) \geq P_H(j)f_{Y|H}(y|j)\} \quad i \neq j,$$

it is easy to see that in region \mathcal{R}_0

$$P_H(0)f_{Y|H}(y|0) \geq P_H(1)f_{Y|H}(y|1)$$

and vice-versa. Thus we can write

$$\begin{aligned}
Pr\{e\} &= P_H(0) \int_{\mathcal{R}_1} f_{Y|H}(y|0)dy + P_H(1) \int_{\mathcal{R}_0} f_{Y|H}(y|1)dy \\
&= \int_{\mathcal{R}_1} \min \{P_H(0)f_{Y|H}(y|0), P_H(1)f_{Y|H}(y|1)\} dy \\
&\quad + \int_{\mathcal{R}_0} \min \{P_H(0)f_{Y|H}(y|0), P_H(1)f_{Y|H}(y|1)\} dy \\
&= \int_{\mathcal{R}_0+\mathcal{R}_1} \min \{P_H(0)f_{Y|H}(y|0), P_H(1)f_{Y|H}(y|1)\} dy \\
&= \int_y \min \{P_H(0)f_{Y|H}(y|0), P_H(1)f_{Y|H}(y|1)\} dy.
\end{aligned}$$

2. To show that for $a, b \geq 0$, $\sqrt{ab} \leq \frac{a+b}{2}$, we proceed as follows. Let $m = (a + b)/2$ be the midpoint of an imaginary segment of the real line that goes from a to b . Let $d = (b - a)/2$ be half the distance between a and b . Writing a and b in terms of m and d we obtain: $ab = (m - d)(m + d) = m^2 - d^2 \leq m^2$ which is the desired result.

Using this and the hint, namely, for $a, b \geq 0$, $\min(a, b) \leq \sqrt{ab}$, we can write

$$\begin{aligned}
Pr\{e\} &= \int_y \min \{P_H(0)f_{Y|H}(y|0), P_H(1)f_{Y|H}(y|1)\} dy \\
&\leq \int_y \sqrt{P_H(0)f_{Y|H}(y|0)P_H(1)f_{Y|H}(y|1)} dy \\
&= \sqrt{P_H(0)P_H(1)} \int_y \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy \\
&\leq \frac{P_H(0) + P_H(1)}{2} \int_y \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy \\
&= \frac{1}{2} \int_y \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy.
\end{aligned}$$

3. In class we upper bounded $Pr\{e|H = 0\}$ and $Pr\{e|H = 1\}$ individually instead of upperbounding the almost final result $Pr\{e\} = P_H(0)Pr\{e|H = 0\} + P_H(1)Pr\{e|H = 1\}$, as we did here. More precisely, what we did in class, written differently, is

$$\begin{aligned}
Pr\{e|H = 0\} &= \int_{\mathcal{R}_1} f_{Y|H}(y|0)dy \\
&= \int_{\mathcal{R}_1} \min \{f_{Y|H}(y|0), f_{Y|H}(y|1)\} dy \\
&\leq \int_{\mathcal{R}_1} \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy \\
&\leq \int_{\mathbb{R}^n} \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy
\end{aligned}$$

The last step, which further loosens the bound, is necessary to find a bound of $Pr\{e|H = 0\}$ that does not depend on \mathcal{R}_1 . This “overbounding” is avoided in (ii) by finding the bound over the whole $Pr\{e\}$.

Problem 5. (*Application of Tight Bhattacharyya Bound*)

(i) Using the *Tight Bhattacharyya Bound*, we get

$$\begin{aligned}
Pr\{e\} &\leq \frac{1}{2} \int_y \sqrt{f_{Y|H}(y|0)f_{Y|H}(y|1)} dy \\
&= \frac{1}{2} \int_y \sqrt{\frac{1}{\sqrt{2\pi\sigma^2}} \exp\left\{-\frac{(y+a)^2}{2\sigma^2}\right\} \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left\{-\frac{(y-a)^2}{2\sigma^2}\right\}} dy \\
&= \frac{1}{2} \int_y \frac{1}{\sqrt{2\pi\sigma^2}} \sqrt{\exp\left\{-\frac{y^2+a^2}{\sigma^2}\right\}} dy \\
&= \frac{1}{2} \exp\left\{-\frac{a^2}{2\sigma^2}\right\} \int_y \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left\{-\frac{y^2}{2\sigma^2}\right\} dy \\
&= \frac{1}{2} \exp\left\{-\frac{a^2}{2\sigma^2}\right\}.
\end{aligned}$$

(ii) The above bound is the same as the one derived in class, which was obtained working specifically with the expression for the Q-function. It is surprising that the *Bhattacharyya Bound*, which applies to arbitrary channels, yields the same result.

Problem 6. (*Bhattacharyya Bound for DMCs*)

1. Inequality (a) follows from the tight *Bhattacharyya Bound*.

Using the definition of DMC, it is straightforward to see that

$$\begin{aligned}
P(\mathbf{y}|\mathbf{s}_0) &= \prod_{i=1}^n P(y_i|s_{0i}) \quad \text{and} \\
P(\mathbf{y}|\mathbf{s}_1) &= \prod_{i=1}^n P(y_i|s_{1i}).
\end{aligned}$$

Inequality (b) follows by substituting the above values in (a).

Equality (c) is obtained by observing that $\sum_{\mathbf{y}}$ is the same as \sum_{y_1, \dots, y_n} (the first one being a vector notation for the sum over all possible y_1, \dots, y_n).

In (c) we see that we want the sum of all possible products. It can also be obtained by summing over each y_i and taking the product of the resulting sum for all y_i . This results in inequality (d). We get equality (e) by writing (d) in a more concise form.

When $s_{0i} = s_{1i}$, $\sqrt{P(y|s_{0i})P(y|s_{1i})} = P(y|s_{0i})$ and thus $\sum_y \sqrt{P(y|s_{0i})P(y|s_{1i})} = \sum_y P(y|s_{0i}) = 1$. This would contribute unity to the product (not really useful!). Thus we are only interested in terms where $s_{0i} \neq s_{1i}$. We form the product of all such sums where $s_{0i} \neq s_{1i}$. We then look out for terms where $s_{0i} = a$ and $s_{1i} = b, a \neq b$ and raise the sum to the appropriate power. (Eg. If we have the product $prpqrpqrr$, we would write it as $p^3q^2r^4$). Hence equality (f).

2. For a binary input channel, we have only two source alphabets $\mathcal{X} = \{a, b\}$. Thus

$$\begin{aligned} Pr\{e\} &\leq z^{n(a,b)} z^{n(b,a)} \\ &= z^{n(a,b)+n(b,a)} \\ &= z^{d_H(\mathbf{s}_0, \mathbf{s}_1)} \end{aligned}$$

3. The value of

- (a) For a binary input Gaussian channel,

$$\begin{aligned} z &= \int_y \sqrt{f_{Y|X}(y|0)f_{Y|X}(y|1)} dy \\ &= \exp\left(-\frac{E}{2\sigma^2}\right) \end{aligned}$$

- (b) For the Binary Symmetric Channel (BSC),

$$\begin{aligned} z &= \sqrt{P(y=0|x=0)P(y=0|x=1)} + \sqrt{P(y=1|x=0)P(y=1|x=1)} \\ &= 2\sqrt{\delta(1-\delta)}. \end{aligned}$$

- (c) For the Binary Erasure Channel (BEC),

$$\begin{aligned} z &= \sqrt{P(y=0|x=0)P(y=0|x=1)} + \sqrt{P(y=E|x=0)P(y=E|x=1)} \\ &\quad + \sqrt{P(y=1|x=0)P(y=1|x=1)} \\ &= 0 + \delta + 0 \\ &= \delta. \end{aligned}$$