## ÉCOLE POLYTECHNIQUE FÉDÉRALE DE LAUSANNE

School of Computer and Communication Sciences

Handout 5	Information Theory and Coding
Homework 3	October 4th, 2010, SG1 – 15:15-17:00

**Problem 1.** Suppose that there are n coins which are either of the same weight or at most one coin has different weight than the others. The coins are to be weighted by a balance. The goal is to find out whether there exists a counterfeit coin and if so, is it heavier or lighter than the others.

- a) For a given positive integer k, find an upper bound on the number of coins, n, for which k usage of the balance is sufficient.
- b) Demonstrate a strategy for n = 12 and k = 3.

**Problem 2.** A biased coin is flipped until both head and tail show up at least once, Suppose that the probability that the outcome of the coin flip is head is p. Let X denote the number of flip required. Find the entropy H(X) in bits.

**Problem 3.** Suppose that  $\mathcal{A}$  is a finite set and P is a probability mass function defined on the set  $\mathcal{A}$ . Let X be a random variable which takes its value from the set  $\mathcal{A}$  with respect to P. Let Y = f(X) where f is some function.

- a) Show that for every function f we have  $H(X) \ge H(Y)$ . Give both intuitive and precise mathematical proof.
- b) Prove that H(X) = H(Y) if and only if f is a one-to-one function.

**Problem 4.** Consider the following code generating for a random variable X which takes on m values  $\{1, 2, ..., m\}$  with probabilities  $p_1 \ge p_2 \ge ... \ge p_m$ . Define  $S_i = \sum_{j=1}^{i-1} p_j$ . Then the codeword for i is the number  $S_i \in [0, 1]$  rounded off to  $\log \lfloor \frac{1}{p_i} \rfloor$  bits (i.e. binary representation of  $S_i$ ). Show that the code constructed by this process is prefix-free and the average length of the code satisfies:  $H(X) \le L < H(X) + 1$ .

**Problem 5.** Find the (a) binary (b) ternary Huffman code for the random variable X with the probability mass function p = (1/12, 1/12, 1/8, 1/8, 1/4, 1/3).